Streams

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Lazy Evaluation in Scheme

Streams are similar to lists, except that the tail of a stream is not evaluated until we asked to do it. This allows streams to be used to represent infinitely long lists.

Lazy evaluation

- In Python, iterators and generators allowed for lazy evaluation

Can represent large or infinite sequences

```
def ints(first):
    while True:
        yield first
        first += 1
```

```
>>> s = ints(1)
>>> next(s)
1
>>> next(s)
2
```

- Scheme doesn't have iterators.

How about a list?

Second argument to cons is always evaluated

```
(define (ints first)
    (cons first (ints (+ first 1))))
```

```
scm> (ints 1)
maximum recursion depth exceeded
```

Streams

Instead of iterators,

Scheme uses streams

```
(define (ints first)
    (cons first
          (ints (+ first 1)))

scm> (ints 1)
maximum recursion depth exceeded
```

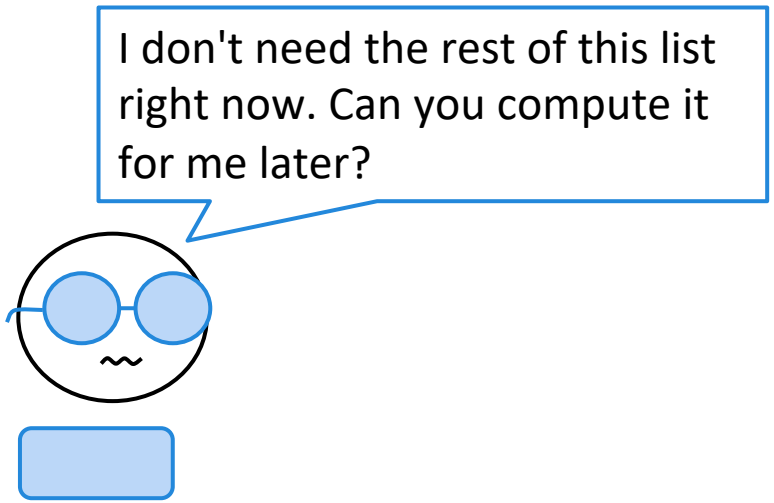
```
(define (ints first)
    (cons-stream first
                 (ints (+ first 1)))

scm> (ints 1)
(1 . #[promise (not forced)])
```

Lazy evaluation, just like iterators in Python

Streams

```
scm> (define s (cons-stream 1 (cons-stream 2 nil)))
s
scm> s
(1 . #[promise (not forced)])
scm> (car s)
1
```

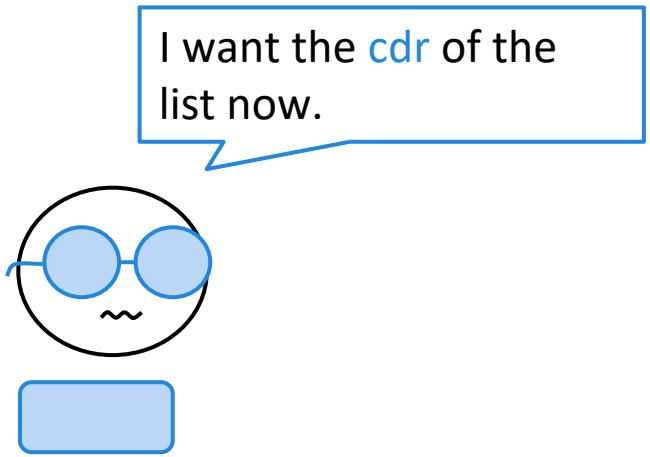
- Stream: (linked) list whose rest is lazily evaluated
  - A promise to compute



Streams

```
scm> (define s (cons-stream 1 (cons-stream 2 nil)))
s
scm> s
(1 . #[promise (not forced)])
scm> (cdr s)
#[promise (not forced)]
```

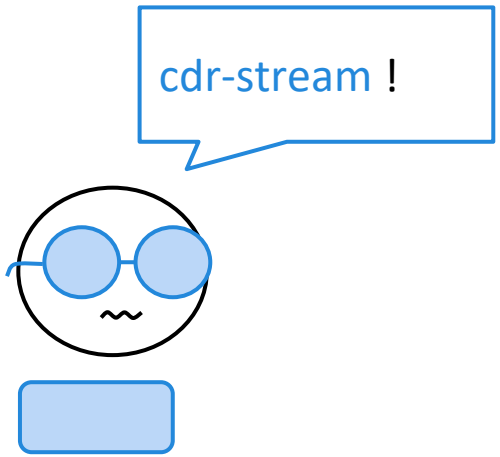
- `cdr` returns the rest of a list
  - For normal lists, the rest is another list
  - The rest of a stream is a promise to compute the list



Streams

```
scm> (define s (cons-stream 1 (cons-stream 2 nil)))
s
scm> s
(1 . #[promise (not forced)])
scm> (cdr-stream s)
(2 . #[promise (not forced)])
scm> (cdr-stream (cdr-stream s))
()
```

- `cdr-stream` forces Scheme to compute the rest



Streams

Remember, a stream is just a regular Scheme pair whose second element is a promise

```
scm> (define s (cons-stream 1 (cons-stream 2 nil)))
s
scm> (cdr s)
#[promise (not forced)]
scm> (cdr-stream s)
(2 . #[promise (not forced)])
scm> (cdr-stream (cdr-stream s))
()
```

Promises: `delay`

- Promise: an object that delays evaluation of an expression
  - The delay special form creates promises

```
scm> (print 5)
5
```

`(print 5)` is immediately evaluated

```
scm> (delay (print 5))
#[promise (not forced)]
```

`(print 5)` is not evaluated yet

Promises: `force`

- The delay special form creates promises
- The force procedure evaluates the expression inside the promise

```
scm> (define x (delay (print 5)))
x
scm> x
#[promise (not forced)]
scm> (force x)
5

scm> x
#[promise (forced)]
```

`(print 5)` is not evaluated yet

Evaluates `(print 5)`

Error, in our interpreter

Promises

`cons-stream` and `cdr-stream` are syntactic sugar.
Achieve the same effect with `delay` and `force`

```
scm> (define s (cons-stream 1 nil)))
s
scm> s
(1 . #[promise (not forced)])
scm> (cdr-stream s)
()
```

```
scm> (define s (cons 1 (delay nil))))
s
scm> s
(1 . #[promise (not forced)])
scm> (force (cdr s))
()
```

Recursively Defined Streams - Constant Stream

Let's start with the constant stream. A constant stream is an infinitely long stream with a number repeated.

```
(define (constant-stream i)
    (cons-stream i (constant-stream i)))
```

```
scm> (define ones (constant-stream 1))
scm> (car ones)
1
scm> (car (cdr-stream ones))
1
```

Check Your Understanding: Natural Number Stream

Let's define the naturals stream which is an infinitely long stream with the natural numbers starting at start.

```
(define (nats start)

  ____________________________)
```

Demo_1

```
scm> (define s (nats 0))
scm> (car s)
0
scm> (car (cdr-stream s))
1
scm> (car (cdr-stream (cdr-stream s)))
2
```

Natural Number Stream

```
(define (nats start)
  (cons-stream start (nats (+ start 1))))
```

Add-Stream and Ints-Stream

Let's write a function that will add two infinite streams together and return a new stream.

```
(define (add-stream s1 s2)
        (cons-stream (+ (car s1) (car s2))
                                 (add-stream (cdr-stream s1)
                                             (cdr-stream s2)))))
```

Let's see it in action! Let's first define the ones stream again.

```
(define ones (cons-stream 1 ones))
```

This is the same as (constant-stream 1).

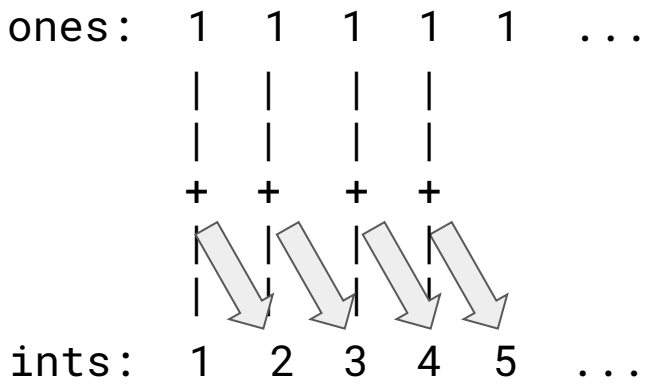
Let's use the ones stream and our new add-stream function to define the ints stream. This is the same as (nats 1). How do we do this?

```
(define ints (cons-stream 1 (add-stream ? ?))
```

Demo_2

Ints-Stream Solution

```
(define ones (cons-stream 1 ones))
(define ints (cons-stream 1 (add-stream ones ints))
```



We can use infinite streams to build other infinite streams. This is the power of lazy evaluation, our current stream stays one step ahead of itself!

Examples: map-stream

- Implement `(map-stream fn s)`:
  - `fn` is a one-argument function
  - `s` is a stream
- Returns a new stream with `fn` applied to elements of `s`

```
(define (map-stream fn s)
  'YOUR-CODE-HERE
)
```

Demo_3

Examples: map-stream

- How would you implement map-list?

```
(define (map-list fn s)
  (if (null? s)
      nil
      (cons (fn (car s))
            (map-list fn (cdr s)))))
```

- How about map-stream?

```
(define (map-stream fn s)
  (if (null? s)
      nil
      (cons-stream (fn (car s))
                   (map-stream fn (cdr-stream s)))))
```

What happens if you change this to cons?

Examples: stream-to-list

- Implement `(stream-to-list s num-elements)`:
  - `s` is a stream
  - `num-elements` is a non-negative integer
- Returns a Scheme list containing the first `num-elements` elements of `s`

```
scm> (stream-to-list (ints 1) 10)
(1 2 3 4 5 6 7 8 9 10)
```

```
(define (stream-to-list s num-elements)
  'YOUR-CODE-HERE
)
```

Demo_4

Examples: stream-to-list

```
(define (stream-to-list s num-elements)
  (if (or (null? s) (= num-elements 0))
      nil
      (cons (car s)
            (stream-to-list (cdr-stream s)
                            (- num-elements 1)))))
)
```